05 - Recursive Algorithm

[KOMS119602] & [KOMS120403]

Design and Analysis of Algorithm (2021/2022)

Dewi Sintiari

Prodi S1 Ilmu Komputer
Universitas Pendidikan Ganesha

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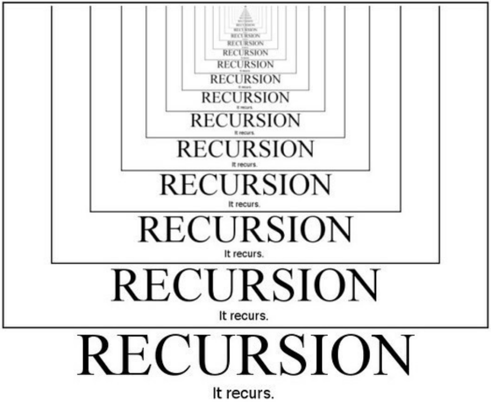


What is recursion or recursive algorithm?

1. The principal of recursive algorithm

A recursive algorithm is an algorithm which calls itself with "smaller (or simpler)" input values, and which obtains the result for the current input by applying simple operations to the returned value for the smaller (or simpler) input.

Characteristics of recursive algorithm:

1. It calls itself recursively
2. It has a base case
3. It must change it state and move towards the base case

A base case is the condition that allows the algorithm to stop recursing: a base case is typically a problem that is small enough to solve directly.

A change of state means that some data that the algorithm is using is modified. Usually the data that represents our problem gets smaller in some way.

Iteration: A function repeats a defined process until a condition fails. This is usually done through a loop, such as a for or while loop with a counter and comparative statement making up the condition that will fail. An infinite loop for iteration occurs when the condition never fails.

Recursion: Instead of executing a specific process within the function, the function calls itself repeatedly until a certain condition is met (this condition being the base case). The base case is explicitly stated to return a specific value when a certain condition is met. An infinite recursive loop occurs when the function does not reduce its input in a way that will converge on the base case.

Simple examples of recursive algorithms

2.1 - Computing factorial (1): Problem statement

$$n! = n \times (n-1) \times (n-2) \times \cdots \times 2 \times 1$$

The formula can be expressed recursively:

$$n! = \begin{cases} n \times (n-1)!, & \text{if } n > 1 \\ 1, & n = 1 \end{cases}$$

2.1 - Computing factorial (2): Pseudocode

Algorithm 1 Factorial of a number

```
procedure FACTORIAL(n)
    if n = 1 then
        return 1
    else
        temp = FACTORIAL(n − 1)
        return n ∗ temp
    end if
end procedure
```

- What is the base case?

Algorithm 2 Factorial of a number

```
procedure FACTORIAL(n)
    if n = 1 then
        return 1
    else
        temp = FACTORIAL(n − 1)
        return n ∗ temp
    end if
end procedure
```

- What is the base case? $n = 1$
- What is the change of states?

2.1 - Computing factorial (2): Pseudocode

Algorithm 3 Factorial of a number

```
procedure FACTORIAL(n)
    if n = 1 then
        return 1
    else
        temp = FACTORIAL(n − 1)
        return n ∗ temp
    end if
end procedure
```

- What is the base case? $n = 1$
- What is the change of states? n decreases
- What is the complexity?

2.1 - Computing factorial (2): Pseudocode

Algorithm 4 Factorial of a number

```
procedure FACTORIAL(n)
    if n = 1 then
        return 1
    else
        temp = FACTORIAL(n − 1)
        return n ∗ temp
    end if
end procedure
```

- What is the base case? $n = 1$
- What is the change of states? n decreases
- What is the complexity? $\mathcal{O}(n)$

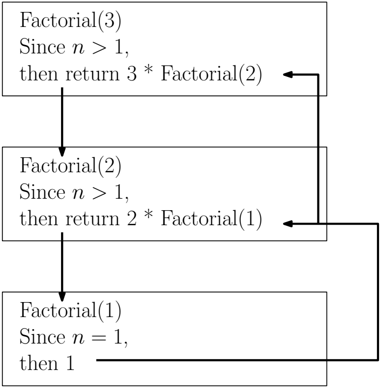


Figure: Illustration of recursive algorithm FACTORIAL with $n = 3$

2.1 - Computing factorial (4): Proving correctness by induction

- Induction base: from line 1, we see that the function works correctly for $n = 1$.
- Hypothesis: suppose the function works correctly when it is called with $n = m$, for some $m \geq 1$.
- Induction step: We prove that it also works when it is called with $n = m + 1$. By the hypothesis, we know the recursive call works correctly for $n = m$ and computes $m!$. Subsequently, it is multiplied by $n = m + 1$, thus computes $(m + 1)!$. And this is the value correctly returned by the program.

To compute the max of n elements for $n > 1$ recursively:

- Compute the max of $n - 1$ elements
- Compare with the last element to find the entire max

2.2 - Finding Maximum Element of an Array (1)

To compute the max of n elements for $n > 1$ recursively:

- Compute the max of $n - 1$ elements
- Compare with the last element to find the entire max

Algorithm 6 Finding maximum of an array

```
procedure MAX(A[0..n − 1], int n)
    if n = 1 then return A[0]
    else
        T = MAX(A, n − 1)
        if T < A[n − 1] then
            return A[n − 1]
        else
            return T
        end if
    end if
end procedure
```

2.2 - Finding Maximum Element of an Array (2)

Task: Compute the following

- Complexity?
- Correctness?

Problem: Given an array of n elements $A[0..n-1]$. We want to compute the sum: $S = \sum_{i=0}^{n-1} A[i]$

Algorithm 7 Sum of an array

```
procedure SUM(A[0..n], int n)
    if n = 1 then return A[0]
    else
        S = SUM(A, n − 1)
        S = S + A[n − 1]
        if T < A[n − 1] then
            return S
        end if
    end if
end procedure
```

Task: Compute the following

- Complexity?
- Correctness?

2.6. Recursive Max, 2nd approach (1)

Problem: Given an array A of n elements, we aim to find an element of maximum value of the array.

Approach:

1. Divide the array into two halves sub-array, namely Left sub-array and Right sub-array.
2. Find the max of each sub-array.
3. Compare the max value of the Left array and the Right array.
4. Return the maximum of the two values.

2.6. Recursive Max, 2nd approach (2)

Algorithm 8 Finding max of an array

procedure $\textsc{FindMax}(A[i..j], n)$ ▷ *i, j are respectively the index of start, end of A*
    if $n = 1$ **then return** $A[S]$
    end if
    $m = \lfloor \frac{i+j}{2} \rfloor$
    $T_1 = \textsc{FindMax}(A[i..m], \lfloor \frac{n}{2} \rfloor)$ ▷ *Recursive call the left sub-array*
    $T_2 = \textsc{FindMax}(A\left[(m+1)..j\right], n - \lfloor \frac{n}{2} \rfloor)$ ▷ *Rec. call right sub-array*
    if $T_1 \geq T_2$ **then return** T_1 ▷ *Compare the two max elements*
    else return T_2
    end if
end procedure

Remark. $\lfloor x \rfloor$ means the largest integer that is $\leq x$;
example: $\lfloor 3.5 \rfloor = 3$

Complexity analysis: Special case when $n = 2^k$

Let $f(n)$: the number of key-comparisons to find the max of an n-array, with $n = 2^k$ for some positive integer k. Hence:

$$f(n) = \begin{cases} 0, & n = 1 \\ 1 + 2f(n/2), & n \geq 2 \end{cases}$$

By repeated substitution:

$$\begin{aligned} f(n) &= 1 + 2f(n/2) \\ &= 1 + 2[1 + 2f(n/4)] = 1 + 2 + 2f(n/4) \\ &= 1 + 2 + 4 + 8f(n/4) \\ &\vdots \\ &= 1 + 2 + 4 + \cdots + 2^{k-1} + 2^k f(n/2^k) \\ &= 1 + 2 + 4 + \cdots + 2^{k-1} \\ &= 2^k - 1/(2-1) = 2^k - 1 \\ &= n - 1 \end{aligned}$$

$f(n)$: the number of key-comparisons to find the max of an n-array, with $n = 2^k$ for some $k \in \mathbb{Z}^+$.

Complexity analysis: For general n

$$f(n) = \begin{cases} 0, & n = 1 \\ f(\lfloor \frac{n}{2} \rfloor) + f(n - \lfloor \frac{n}{2} \rfloor) + 1, & n \geq 2 \end{cases}$$

Prove that:

By induction, we obtain $f(n) = n - 1$. How?

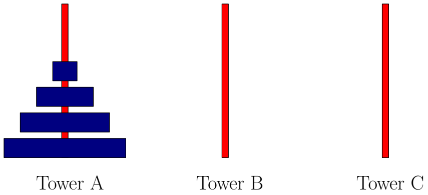
Tower of Hanoi

3. Tower of Hanoi problem (1): Problem statement

Problem: there are three towers A, B, and C. Initially, there are n disks of varying sizes stacked on tower A, ordered by their size, with the largest disk on the bottom and the smallest one on the top. The objective is to move all discs to the 2nd tower by keeping their order.

- Only one disk may be moved at a time in a restricted manner, from the top of one tower to the top of another tower.
- A larger disk must never be placed on top of a smaller disk.

Check `https://www.mathsisfun.com/games/towerofhanoi.html` for an illustration of the problem



3. Tower of Hanoi problem (2): Illustration

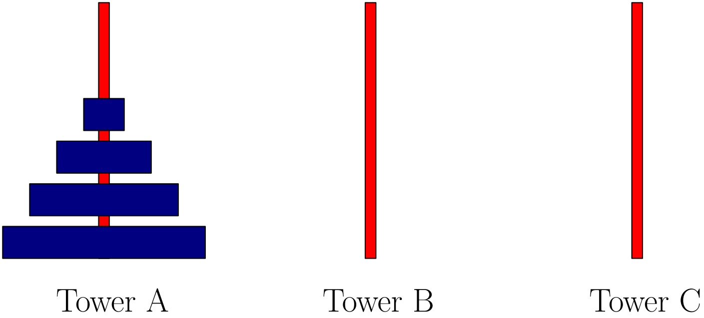


Figure: Initial configuration with 4 disks on Tower A

3. Tower of Hanoi problem (2): Illustration

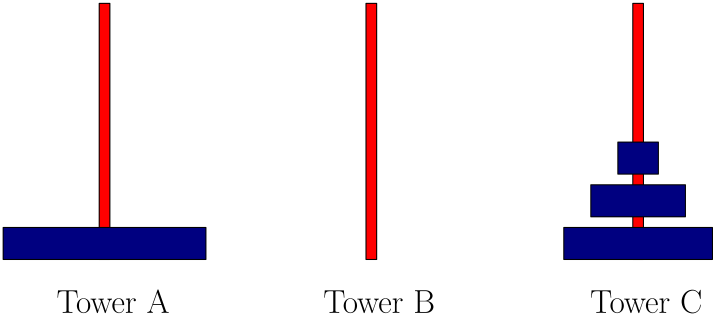


Figure: After recursively moving the top 3 disks from Tower A to Tower C

3. Tower of Hanoi problem (2): Illustration

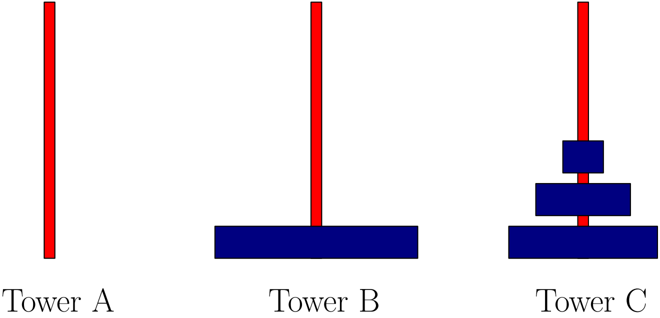


Figure: After moving the bottom disk from Tower A to Tower B

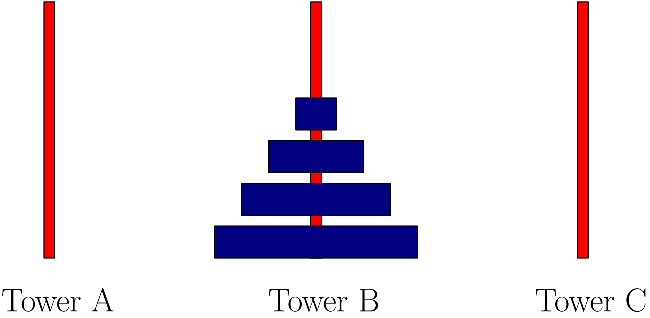


Figure: After recursively moving back 3 disks from Tower C to Tower B.

3. Tower of Hanoi problem (3): Pseudocode

Task: Write the pseudocode of the Tower of Hanoi problem!

Task: Write the pseudocode of the Tower of Hanoi problem!

Algorithm 10 Tower of Hanoi

```
procedure TOWERS(A, B, C, n)
    if n = 1 then
        MOVEONE(A, B)
        return
    end if
    TOWERS(A, C, B, n − 1)
    MOVEONE(A, B)
    TOWERS(C, B, A, n − 1)
end procedure
```

- $\text{TOWERS}(A, B, C, n)$: move n disks from A to B, using A, B, C
- $\text{MOVEONE}(A, B)$: move one disk from A to B

3. Tower of Hanoi problem (4): Proof of correctness

Correctness: by induction

- Base case: For $n = 1$, a single move is made from A to B. So the algorithm works correctly for $n = 1$.
- For any $n \geq 2$, suppose the algorithm works correctly for $n - 1$.
- Then, by the hypothesis, the recursive call of line 6 works correctly and moves the top $n - 1$ disks to C, leaving the bottom disk on tower A.
- The next step, line 7, moves the bottom disk to B.
- Finally, the recursive call of line 8 works correctly by the hypothesis and moves back $n - 1$ disks from C to B.
- Thus, the entire algorithm works correctly for n.

Recurrence equation to analyze time complexity

Let $f(n)$: the number of single moves to solve the problem for n disks

Hence we have the following relation:

$$f(n) = \begin{cases} 1, & \text{if } n = 1 \\ 1 + 2f(n-1), & n \geq 2 \end{cases}$$

Remark. The above formula is known as recursive formula; read this page or watch this video to get an overview.

To obtain the explicit formula of $f(n)$, we have to solve the recurrence equation for $f(n)$.

Method 1: Repeated substitution

$$
\begin{aligned}
f(n) &= 1 + 2 \cdot f(n-1) \\
&= 1 + 2 + 4 \cdot f(n-2) \\
&= 1 + 2 + 4 + 8 \cdot f(n-3) \\
&= \cdots \\
&= 1 + 2 + 2^2 + \cdots + 2^{n-1} \cdot f(1)
\end{aligned}
$$

Substituting the base case $f(1) = 1$ and by the geometric sum formula (click here to check the formula), we obtain:

$$f(n) = \frac{2^n - 1}{2 - 1} = 2^n - 1$$

Method 2: Guess the solution and prove by induction

Suppose our guess is "$f(n)$ is exponential"

Guess: $f(n) = a \cdot 2^n + b$

Inductive proof:

- Induction base: $n = 1$
  - $f(1) = 1$ (from the reccurence)
  - $f(1) = 2a + b$ (from the solution form)

  So we have $2a + b = 1$
- Induction: Suppose that the solution is correct for some $n \geq 1$:

$$f(n) = a \cdot 2^n + b$$

  Then the solution must hold for $n + 1$, that is:

$$f(n+1) = a \cdot 2^{n+1} + b$$

- From the recurrence relation, we obtain:

$$\begin{aligned} f(n+1) &= 2f(n)+1 \\ &= 2(a\cdot 2^n + b)+1 \\ &= a\cdot 2^{n+1} + (2b+1) \end{aligned}$$

- From the two equations, we obtain:

$$a\cdot 2^{n+1} + b = a\cdot 2^{n+1} + (2b+1) \Leftrightarrow 2b+1 = b \Leftrightarrow b = -1$$

  Hence, $2a+b=1 \Leftrightarrow a=1$. So, $b=-1$.

- Hence, $f(n) = a\cdot 2^n + b = 2^n - 1$.

Binary Search

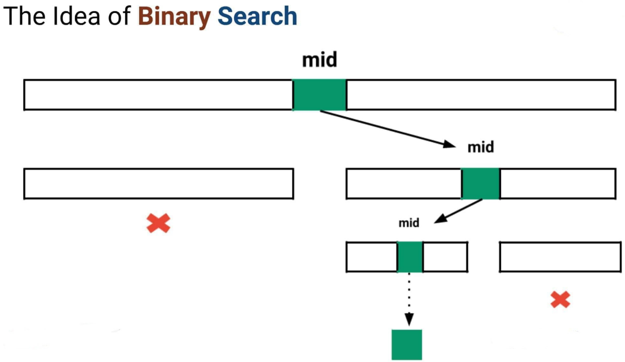
4. Binary search algorithm (1): Principal

Problem: Given a *sorted* array $A[0..n-1]$ and a search key KEY. The algorithm does the following:

- If $KEY = A[m]$, then return m
- If $KEY < A[m]$, then recursively search the left half of the array
- If $KEY > A[m]$, then recursively search the right half of the array

In each step, the size of the search *is reduced by half*.

4. Binary search algorithm (2): Diagram



source: `https://www.enjoyalgorithms.com/blog/binary-search-algorithm`

Algorithm 11 Binary search algorithm

procedure $\textsc{BinSearch}(A, i, j, KEY)$
  if $i > j$ **then**
    return -1 ▷ *Base case is reached but KEY is not found*
  end if
  $m = \lfloor \frac{i+j}{2} \rfloor$ ▷ *Choose the pivot*
  if $KEY = A[m]$ **then**
    return m ▷ *KEY is found at index m*
  else
    if $KEY < A[m]$ **then** ▷ *The KEY is located on the Left sub-array*
      return $\textsc{BinSearch}(A, i, m-1, KEY)$ ▷ *Rec-call left part*
    else
      return $\textsc{BinSearch}(A, m+1, j, KEY)$ ▷ *Rec-call right part*
    end if
  end if
end procedure

4. Binary search algorithm (4): Time complexity analysis

Let $f(n)$ be the number of comparisons.

Complexity analysis: special case when $n = 2^k$

$$f(n) = \begin{cases} 1, & n = 1 \\ 1 + f(n/2), & n \geq 2 \end{cases}$$

By repeated substitution:

$$\begin{aligned} f(n) &= 1 + f(n/2) \\ &= 1 + 1 + f(n/4) \\ &= 1 + 1 + 1 + f(n/8) \\ &\vdots \\ &= k + f(n/2^k) \\ &= k + f(1) \\ &= k + 1 \\ &= \log n + 1 \end{aligned}$$

Complexity analysis: general case n

$$f(n) = \begin{cases} 1, & n = 1 \\ 1 + f(\lfloor \frac{n}{2} \rfloor), & n \geq 2 \end{cases}$$

By induction, we obtain $f(n) = \lfloor \log n \rfloor + 1$

Exercise: show it!

4. Binary search algorithm (6): Inductive proof for TC analysis

- Induction base: $n = 1$:
  From the recurrence, $f(1) = 1$, and the claimed solution $f(1) = \lfloor \log 1 \rfloor + 1 = 1$. Correct.

- Inductive proof: Suppose that the formula is correct for all smaller values.

$$f(m) = \lfloor \log m \rfloor, \quad \forall m < n$$

Every integer n can be expressed as (for some integer k):

$$2^{k-1} \leq \lfloor n/2 \rfloor < 2^k$$

So, $\lfloor \log \lfloor n/2 \rfloor \rfloor = k - 1$.

By the recursive function:

$$f(\lfloor n/2 \rfloor) = \lfloor \log \lfloor n/2 \rfloor \rfloor + 1 = (k-1) + 1 = k = \lfloor \log n \rfloor$$

Then:

$$f(n) = f(\lfloor n/2 \rfloor) + 1 = k + 1 = \lfloor \log n \rfloor + 1$$

A more advanced example: Recursive powering

5. Recursive powering (1): Problem statement

Problem: Given X and an integer n. We want to compute X^n.

Algorithm 12 Recursive powering (*brute force*)

```
procedure POWER1(X, n)
    T = X
    for i = 2 to n do
        T = T * X
    end for
end procedure
```

Complexity: $\mathcal{O}(n)$. Why?

Idea: $X^{16} = ((((X^2)^2)^2)^2$

Given $n = 2^k$, we can do repeated squaring.

Algorithm 13 Recursive powering (*special case*)

```
procedure POWER2(X, n = 2^k)
    T = X
    for i = 2 to k do
        T = T * T
    end for
end procedure
```

Complexity: $\mathcal{O}(\log n)$. Why?

Generalize the case for n: Computing X^n for $n \in \mathbb{Z}^+$

- Compute $X^2 = X * X$
- Compute $X^3 = X^2 * X$
- Compute $X^6 = X^3 * X^3$
- Compute $X^{12} = X^6 * X^6$
- Compute $X^{13} = X^{12} * X$

5. Recursive powering (4): Approach

Basic idea: Divide n by 2, $n = n/2 + n/2$. So

$$X^n = X^{(n/2+n/2)} = X^{n/2} \cdot X^{n/2}$$

The problem is $n/2$ is not always an integer. So we have to apply a little modification:

- For $n = 0$, then $X^n = 1$
- For $n > 0$, then:
  - If n is *even*, then $X^n = X^{n/2} \cdot X^{n/2}$
  - If n is *odd*, then $X^n = X^{\lfloor n/2 \rfloor} \cdot X^{\lfloor n/2 \rfloor} \cdot X$

5. Recursive powering (5): Pseudocode

Algorithm 14 Recursive powering

```
procedure POWER3(X, n)
    if n = 1 then
        return X
    end if
    T = POWER(X, ⌊n/2⌋)
    T = T * T
    if n mod 2 = 1 then
        T = T * X
        return T
    end if
end procedure
```

Complexity: ?

Example: Computing 3^{16}

$$
\begin{aligned}
3^{16} = 3^8 \cdot 3^8 &= (3^8)^2 \\
&= ((3^4)^2)^2 \\
&= (((3^2)^2)^2)^2 \\
&= ((((3^1)^2)^2)^2)^2 \\
&= ((((3^0) \cdot 3)^2)^2)^2)^2 \\
&= (((1 \cdot 3)^2)^2)^2)^2 \\
&= (((3)^2)^2)^2)^2 \\
&= (((9)^2)^2)^2 \\
&= ((81)^2)^2 \\
&= (6561)^2 \\
&= 43,046,721
\end{aligned}
$$

5. Recursive powering (7): Correctness (*informal proof*)

Algorithm 14 Power by multiplications

```
procedure POWER3(X, n)
    if n = 1 then
        return X
    end if
    T = POWER(X, ⌊n/2⌋)
    T = T * T
    if n mod 2 = 1 then
        T = T * X
        return T
    end if
end procedure
```

Let $n = 2m + r$, where $r \in \{0, 1\}$.

- The algorithm first makes a recursive call to compute $T = X^m$.
- Then it squares T to get $T = X^{2m}$. If $r = 0$, this is returned.
- Otherwise, when $r = 1$, the algorithm multiplies T by X, to result in $T = X^{2m+1}$.

Let $f(n)$: the worst-case number of multiplication steps to compute X^n.

- The recursive call takes $f(\lfloor \frac{n}{2} \rfloor)$ multiplications.
- Then it is followed by one more multiplication. In the worst case, when n is odd, one additional multiplication is performed.

Hence,

$$f(n) = \begin{cases} 0, & \text{if } n = 1 \\ f(\lfloor \frac{n}{2} \rfloor) + 2, & \text{if } n \geq 2, n \text{ odd} \\ f(\lfloor \frac{n}{2} \rfloor) + 1, & \text{if } n \geq 2, n \text{ even} \end{cases}$$

Show that $f(n) = 2\lfloor \log n \rfloor$.

5. Recursive powering (8): Time complexity analysis

$$f(n) = \begin{cases} 0, & \text{if } n = 1 \\ f(\lfloor \frac{n}{2} \rfloor) + 2, & \text{if } n \geq 2, n \text{ odd} \\ f(\lfloor \frac{n}{2} \rfloor) + 1, & \text{if } n \geq 2, n \text{ even} \end{cases}$$

The last two cases have small difference. So we can approximate the function above with the following function to simplify the computation:

$$f(n) = \begin{cases} 0, & \text{if } n = 1 \\ f(\lfloor \frac{n}{2} \rfloor) + 2, & \text{if } n \geq 2 \end{cases}$$

5. Recursive powering (9): Inductive proof

Show that $f(n) = 2\lfloor \log n \rfloor$.

- Induction base ($n = 1$): From the recurrence, $f(1) = 0$, and from the formula, $f(1) = 2\lfloor \log 1 \rfloor = 0$. Correct.
- Inductive proof: Suppose that the formula is correct for all smaller values.

$$f(m) = 2\lfloor \log m \rfloor, \quad \forall m < n$$

Every integer n can be expressed as (for some integer k):

$$2^k \leq n < 2^{k+1}$$

So, $\lfloor \log n \rfloor = k$, and $\lfloor \frac{\log n}{2} \rfloor = k - 1$. By the recursive function:

$$f(n) = f(\lfloor \frac{n}{2} \rfloor) + 2 = 2(k-1) + 2 = 2k = 2\lfloor \log n \rfloor$$

Remark. This gives a better complexity than the brute force approach ($\mathcal{O}(n)$).

Redundancy in recursive algorithm

Redundancy (1): Recursive powering

Algorithm 14 Power by multiplications

procedure POWER3(X, n)
    if $n = 1$ **then**
        return X
    end if
    $T = \textsc{Power}(X, \lfloor \frac{n}{2} \rfloor)$
    $T = T * T$
    if $n \mod 2 = 1$ **then**
        $T = T * X$
        return T
    end if
end procedure

Is it necessary to store $\textsc{Power}(X, \lfloor \frac{n}{2} \rfloor)$ in some variable T?

Redundancy (2): Recursive powering

Algorithm 15 Recursive powering

```
procedure POWER4(X, n)
    if n = 1 then
        return X
    end if
    return POWER(X, ⌊n/2⌋) * POWER(X, ⌊n/2⌋)
end procedure
```

- Is the algorithm correct? *What is the complexity?*

Redundancy (3): Recursive powering

The algorithm is correct.

The number of recursive calls:

$$f(n) = \begin{cases} 0, & \text{if } n = 1 \\ f(\lfloor \frac{n}{2} \rfloor) + f(\lfloor \frac{n}{2} \rfloor) + 1, & \text{if } n \geq 2 \end{cases}$$

By induction, we can prove that $f(n) = n - 1$ (asymptotically worse than the previous algorithm).

What can you conclude?

Redundancy (3): Recursive powering

The algorithm is correct.

The number of recursive calls:

$$f(n) = \begin{cases} 0, & \text{if } n = 1 \\ f(\lfloor \frac{n}{2} \rfloor) + f(\lfloor \frac{n}{2} \rfloor) + 1, & \text{if } n \geq 2 \end{cases}$$

By induction, we can prove that $f(n) = n - 1$ (asymptotically worse than the previous algorithm).

What can you conclude?
POWER4 is also not efficient, because we make two recursive calls for the same function $f(\lfloor \frac{n}{2} \rfloor)$

The Fibonacci sequence is defined as follows.

$$F(n) = \begin{cases} 1, & n = 1 \\ 1, & n = 2 \\ F_{n-1} + F_{n-2}, & n \geq 3 \end{cases}$$

Fibonacci sequence: $1, 2, 3, 5, 8, 13, 21, \ldots$

Redundancy (4): Fibonacci sequence

The Fibonacci sequence is defined as follows.

$$F(n) = \begin{cases} 1, & n = 1 \\ 1, & n = 2 \\ F_{n-1} + F_{n-2}, & n \geq 3 \end{cases}$$

Fibonacci sequence: $1, 2, 3, 5, 8, 13, 21, \ldots$

Build an algorithm to compute the Fibonacci sequence!

- With naive algorithm (brute force), we can reach complexity $\mathcal{O}(n)$. How?

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Fibonacci sequence: $1, 2, 3, 5, 8, 13, 21, \ldots$

Build an algorithm to compute the Fibonacci sequence!

- With naive algorithm (brute force), we can reach complexity $\mathcal{O}(n)$. How?

  By looping (iterative method); we add the number one by one.

Redundancy (4): Fibonacci sequence

The Fibonacci sequence is defined as follows.

$$F(n) = \begin{cases} 1, & n = 1 \\ 1, & n = 2 \\ F_{n-1} + F_{n-2}, & n \geq 3 \end{cases}$$

Fibonacci sequence: $1, 2, 3, 5, 8, 13, 21, \ldots$

Build an algorithm to compute the Fibonacci sequence!

- With naive algorithm (brute force), we can reach complexity $\mathcal{O}(n)$. How?

  By looping (iterative method); we add the number one by one.

- Create a recursive algorithm!

Redundancy (5): Fibonacci sequence

Algorithm 16 Fibonacci sequence

```
procedure Fib(n)
    if n ≤ 2 then return 1
    end if
    return (Fib(n − 1) + Fib(n − 2))
end procedure
```

Redundancy (5): Fibonacci sequence

Algorithm 17 Fibonacci sequence

```
procedure Fib(n)
    if n ≤ 2 then  return 1
    end if
    return (Fib(n − 1) + Fib(n − 2))
end procedure
```

This program makes recursive calls with a great deal of overlapping computations, causing a huge inefficiency.

Redundancy (5): Fibonacci sequence

Algorithm 18 Fibonacci sequence

```
procedure FIB(n)
    if n ≤ 2 then return 1
    end if
    return (FIB(n − 1) + FIB(n − 2))
end procedure
```

This program makes recursive calls with a great deal of overlapping computations, causing a huge inefficiency.

Complexity:

$$T(n) = \begin{cases} 0, & n = 1 \\ 0, & n = 2 \\ T(n-1) + T(n-2) + 1, & n \geq 3 \end{cases}$$

Prove that: the explicit function: $T(n) \geq (1.618)^{n-2}$.

Advantages and drawbacks of recursive algorithm (1)

Advantages

- Recursion adds clarity and reduces the time needed to write and debug code (since it reduce the length of code).
- To solve such problems which are naturally recursive such as tower of Hanoi.
- Recursion can reduce time complexity (*sometimes counter-intuitive*).
- Reduce unnecessary calling of function.
- Extremely useful when applying the same solution.

Advantages and drawbacks of recursive algorithm (2)

Drawbacks

- Recursive functions are generally slower than non-recursive function.
- It may require a lot of memory space to hold intermediate results on the system stacks.
- Hard to analyze or understand the code.
- It is not more efficient in terms of space and time complexity (can be slow).
- The computer may run out of memory if the recursive calls are not properly checked.

Sum up...

What have we learned today?

1. Reviewing brute force approach
2. Understanding the principal of recursive approach
3. Some examples of recursive algorithms
4. Recurrence equation to analyze time complexity
5. Redundancy in recursion: be careful when writing the pseudocode
6. Binary search algorithm