TUTORIAL XI

1 Singleton Bound

For every $(n, k, d)_q$-code, show that $k \leq n - d + 1$.

2 Weights of Codewords

Let C be an $[n, k, d]$-linear code over $\mathbb{F}_q$. Prove the following.

1. For $q = 2$, either all the codewords have even weight or exactly half have even weight and the rest have odd weight.
2. For any q, either all the codewords begin with 0 or exactly a fraction $1/q$ of the codewords begin with 0. In general, for a given position $1 \leq i \leq n$, either all codewords contain 0 at the i-th position or each $\alpha \in \mathbb{F}_q$ appears at the i-th position of exactly $1/q$ of the codewords in C.
3. The following inequality holds for the minimum distance d of C.

$$d \leq \frac{n(q-1)q^{k-1}}{q^k - 1}$$

3 Codes Achieving the Gilbert-Varshamov Bound

The purpose of this exercise is to use the probabilistic method to show that a random linear code lies on the Gilbert-Varshamov bound, with high probability.

1. Given a non-zero vector $\mathbf{m} \in \mathbb{F}_q^k$ and a uniformly random $k \times n$ matrix $\mathbf{G}$ over $\mathbb{F}_q$, show that the vector $\mathbf{mG}$ is uniformly distributed over $\mathbb{F}_q^n$.
2. Let $k = (1 - H_q(\delta) - \varepsilon)n$, with $\delta = d/n$. Show that there exists a $k \times n$ matrix $\mathbf{G}$ such that

$$\text{for every } \mathbf{m} \in \mathbb{F}_q^k \setminus \{\mathbf{0}\}, wt(\mathbf{mG}) \geq d$$

   where $wt(\mathbf{m})$ is the Hamming weight of the vector $\mathbf{m}$.
3. Show that $\mathbf{G}$ has full rank (i.e., it has dimension at least $k = (1 - H_q(\delta) - \varepsilon)n$)